

**Handout 2**

Let  $S$  (and sometimes  $X$ ) be sets of outcomes with typical elements  $x, y, z$ , or, sometimes,  $x_1, x_2, \dots, x_n$ .

**Preference relations**

**Def<sup>n</sup>:**  $x \succ_i y$  means that person  $i$  *strictly prefers* outcome  $x$  to outcome  $y$ .

**Def<sup>n</sup>:**  $x \succeq_i y$  (“ $x$  is weakly preferred to  $y$  by  $i$ ”) if it is not the case that  $y \succ_i x$ .

Or equivalently, not  $y \succ x \Leftrightarrow x \succeq y$ . The  $\Leftrightarrow$  means or reads “if and only if.”

**Def<sup>n</sup>:**  $x \sim_i y$  ( $i$  is “indifferent” between outcomes  $x$  and  $y$ ) if it is not the case that  $x \succ_i y$  and it is not the case that  $y \succ_i x$ .

An implication: not  $x \succ y \Leftrightarrow y \succeq x$ , and not  $y \succ x \Leftrightarrow x \succeq y$ . So we could also write  $x \sim_i y$  if  $x \succeq y$  and  $y \succeq x$ .

**Axioms about preferences**

A1: (completeness) For any  $x, y \in S$ , either  $x \succeq y$  or  $y \succeq x$ .

A2: (Transitivity of  $\succeq$ ) For all  $x, y, z \in S$ , if  $x \succeq y$  and  $y \succeq z$ , then  $x \succeq z$ .

**Numerical representations of preferences (ordinal utility)**

**Defn:**  $u : S \rightarrow R$  is a numerical representation of  $\succeq$  on  $S$  if

$$u(x) \geq u(y) \Leftrightarrow x \succeq y.$$

**Equivalent ordinal utility functions**

**Claim:** If  $u(x)$  represents a person’s preferences  $\succeq$  on  $S$  in the sense just defined, then so does any function  $v(x)$  that is a monotonic (i.e., order preserving) transformation of  $u(x)$ .

**Conditions for an ordinal utility representation to exist**

**Thm:** For  $S$  finite, there exists a numerical representation  $u$  of  $\succeq$  iff  $\succeq$  satisfies A1 and A2.

Sketch of Proof:

Two parts: (1) If there exists a  $u$  that is a n.r. then A1 and A2 must be satisfied.

Since  $u(x)$  and  $u(y)$  are numbers, it is the case that either  $u(x) \geq u(y)$  or  $u(y) \geq u(x)$ . So, from the defn of a numerical representation,  $u(x) \geq u(y)$  implies  $x \succeq y$  and  $u(y) \geq u(x)$  implies  $y \succeq x$ , so completeness of  $\succeq$  must obtain.

Similarly for transitivity.

Second part: If a dmkr satisfies A1 and A2, then there exists a  $u$  that is a numerical representation of  $\succeq$ .

To prove this, show *by construction*, an important proof technique. Use the axioms to show how to construct a  $u$  that works.

First step: Show that A1 and A2 imply that we can order the outcomes in  $S$  as  $x_i \succeq x_j \succeq x_k \dots$ , using  $\sim$  when appropriate. Completeness implies that dmkr can say " $x \succeq y$ " or " $y \succeq x$ " to every pairing, and if it were not possible to order the outcomes like this then transitivity (A2) must be violated for some triple of outcomes.

Second step: Assign numbers to each outcome in descending order, assigning the same number whenever  $x_i \sim x_j$ .

Third step: Show that by construction this u fn is a numerical representation of dmkr's preferences in the sense defined above.

### **Axioms, etc., using the strict preference relation**

A1': (asymmetry) For all  $x, y \in S$ ,  $x \succ y$  implies not  $y \succ x$ .

Claim: A1 is true iff A1' is true.

Proof: (will use the technique of proof by contradiction, in which one assumes the claim is false and then shows that this generates a contradiction).

First, suppose A1 is true and A1' is false. This would mean that

A1 true: for all  $x, y \in S$ , either  $x \succeq y$  or  $y \succeq x$ , AND  
A1' false: for some  $x, y \in S$ ,  $x \succ y$  and  $y \succ x$ .

Remember that  $\succeq$  and  $\succ$  are defined in terms of each other. Using the fact that  $x \succeq y$  implies not  $y \succ x$  by definition, we can rewrite this

A1 true: for all  $x, y \in S$ , either not  $y \succ x$  or not  $x \succ y$ , AND  
A1' false: for some  $x, y \in S$ ,  $x \succ y$  and  $y \succ x$ .

But notice now that A1 and A1' cannot both be true. So we have proven that if A1 is true then so must A1' be also.

What about the other direction? Suppose next that A1' is true and A1 is false. Try to do it.

A2'': (Transitivity of  $\succ$ ) For all  $x, y, z \in S$  such that  $x \succ y$  and  $y \succ z$ , it is true that  $x \succ z$ .

A2': (negative transitivity) For any  $x, y, z \in S$ , if  $x \succ y$ , then either  $x \succ z$  or  $z \succ y$  or both.

Claim: A2 is true if and only if A2' is true.

Claim: A2' implies A2'', but A2'' does not imply A2'.

### Lotteries, or gambles

**Def<sup>n</sup>**: Let  $X = \{x_1, x_2, \dots, x_m\}$  be a set of outcomes.

$$\Delta(X) = \{(p_1, p_2, \dots, p_n) : \sum_{i=1}^n p_i = 1 \text{ and } p_i \geq 0 \text{ for all } i\}$$

is the set of all probability distributions on the set  $X$ , which we will also call the set of *lotteries* or *gambles* on  $X$ .

### The “natural principle”

A lottery that puts more weight (in terms of probability) on *good* outcomes (from the point of view of the decision maker), and less weight on *bad* outcomes is preferable to a lottery that puts more weight on the bad outcomes and less on the good ones.

### An “ordinal utility” representation of preferences over lotteries

If a person's preferences over lotteries in  $\Delta(X)$  satisfy A1, A2, and A3 (continuity, not given yet) then there is a utility function  $u : \Delta(X) \rightarrow \mathbb{R}$  that represents these preferences in the sense that for all  $p, q \in \Delta(X)$ ,  $p \succ q$  iff  $u(p) > u(q)$ .

To observe: Preferences that can be represented in this way need not satisfy the natural principle.

### Expected value

**Def<sup>n</sup>**: The *expected value* of a lottery  $p \in \Delta(X)$  is defined as

$$EV(p) = \sum_{i=1}^n p_i x_i.$$

Note that  $EV(p)$  is defined only if the set of outcomes is a set of numbers.

### An expected value representation of preferences over lotteries

**Def<sup>1</sup>:** A person's preferences are represented by the expected value function iff for  $p, q \in \Delta(X)$ ,  $p \succ q \Leftrightarrow EV(p) > EV(q)$ .

### The expected utility representation of preferences on $\Delta(X)$

**Def<sup>2</sup>:** A person's preferences over the lotteries in  $\Delta(X)$  (where  $X$  is a set of outcomes) can be represented by an *expected utility function* if there exists a function  $u : X \rightarrow \mathbb{R}$  such that for any  $p, q \in \Delta(X)$ ,

$$p \succ q \Leftrightarrow EU(p) \equiv \sum_{i=1}^n p_i u(x_i) > \sum_{i=1}^n q_i u(x_i) \equiv EU(q).$$

### Equivalent expected utility functions

Claim: If  $u(x)$  is an expected utility function that represents a person's preferences over gambles in  $\Delta(X)$ , then so is any  $v(x) = au(x) + b$ , where  $a > 0$  and  $a, b \in \mathbb{R}$ .

Proof: Suppose  $u(x)$  represents  $\succ$  on  $\Delta(X)$  in the sense that for any  $p, q \in \Delta(X)$ ,  $p \succ q \Leftrightarrow EU(p) > EU(q)$ . Thus,

$$\begin{aligned} \sum p_i u(x_i) &> \sum q_i u(x_i) \text{ (recall what } p_i, q_i \text{ are)} \\ a \sum p_i u(x_i) &> a \sum q_i u(x_i) \\ a(\sum p_i u(x_i)) + b &> a(\sum q_i u(x_i)) + b \\ \sum p_i (au(x_i) + b) &> \sum q_i (au(x_i) + b) \\ \sum p_i v(x_i) &> \sum q_i v(x_i). \end{aligned}$$